

Packing Nearly Optimal Ramsey $R(3, t)$ Graphs

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ABSTRACT/SUMMARY

- Previous work: Kim/Bohman construct one nearly optimal $R(3, t)$ graph via (semi-random) triangle-free process
- Our contribution: We approximately decompose K_n into nearly optimal $R(3, t)$ graphs via iterative semi-random Δ -free process
- Our application: We solve Ramsey-Conjecture by Fox et al. (determine order of Ramsey-type parameter of Burr-Erdős-Lovász from 1976)

Kim (1995), Bohman (2008): one nearly optimal $R(3, t)$ graph

They both find an n -vertex graph $G \subseteq K_n$ such that G is Δ -free with independence number $\alpha(G) \leq C\sqrt{n \log n}$

- Kim/Bohman use (semi-random) Δ -free process
- Δ -free process
In each step: add one random edge that does not close a Δ
- Semi-random variation of Δ -free process (Rödl nibble type)
In each step: add “many” random-like edges that do not close a Δ

G., Warnke (2020): almost packing of nearly optimal $R(3, t)$ graphs

Given $\epsilon > 0$, we find edge-disjoint $(G_i)_{i \in \mathcal{I}}$ with $G_i \subseteq K_n$ such that
(a) each n -vertex graph G_i is Δ -free with $\alpha(G_i) \leq C_\epsilon \sqrt{n \log n}$
(b) $|\bigcup_{i \in \mathcal{I}} E(G_i)| \geq (1 - \epsilon) \binom{n}{2}$

- Via simple polynomial-time randomized algorithm:
Sequentially choose G_i via semi-random variation of Δ -free process
- Main technical challenge:
Controlling errors over $\Theta(\sqrt{n/\log n})$ iterations of the process

PROOF: HIGH-LEVEL STRATEGY/IDEAS

Main technical result: find random-like Δ -free subgraph $G \subseteq H$

Let $p := \sqrt{\beta(\log n)/n}$ and $s := C_\epsilon \sqrt{n \log n}$. If $H \subseteq K_n$ is such that

$$e_H(A, B) \geq \epsilon |A||B|$$

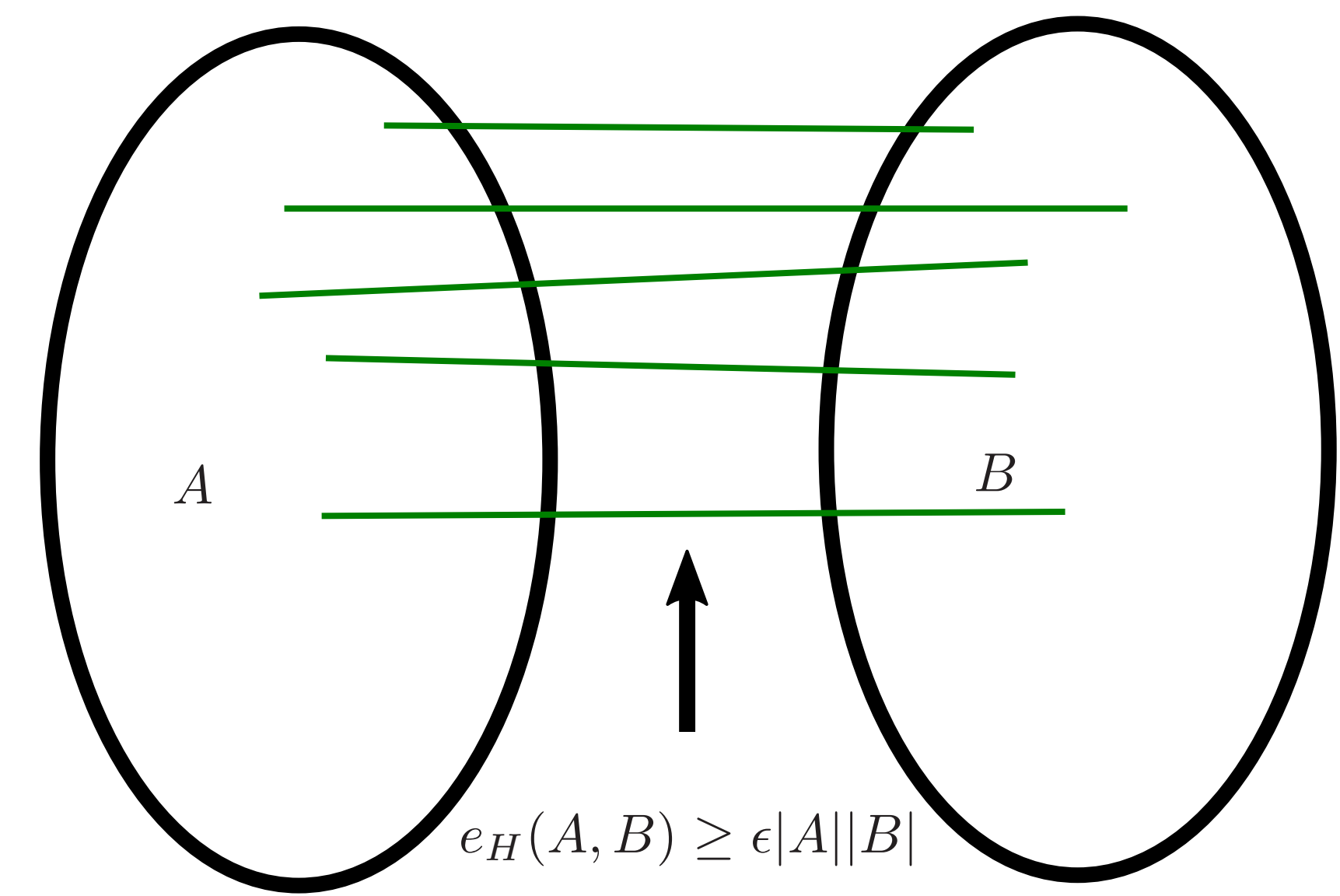
for all disjoint sets A, B of size s , then we can find Δ -free $G \subseteq H$ with

$$e_G(A, B) = (1 \pm \delta) p e_H(A, B)$$

for all disjoint sets A, B of size s .

Proof based on semi-random variation of Δ -free process:

- Crucial that we do not require degree/codegree regularity of H
- “Self-stabilization” mechanism built into process (to control errors)
- One-sided error estimates (only track large-sets two-sided)
- Tools: Bounded-Differences-Ineq., and Upper-Tail-Ineq. of Warnke



Algorithm for packing result: (maintain $e_{H_i}(A, B)$ bounds inductively)

- Start with $H_0 = K_n$
- Sequentially choose $G_i \subseteq H_i$ and set $H_{i+1} = H_i \setminus G_i$, noting
 $e_{H_i}(A, B) = (1 - (1 \pm \delta)p)^i |A||B|$
- Stop when $e_{H_i}(A, B) \approx \epsilon |A||B|$ holds

APPLICATION: SOLVE RAMSEY-CONJECTURE OF FOX ET AL.

Ramsey-Definitions for $r \geq 2$ colors

- $G \rightarrow (\mathbf{K}_k)_r$ if any r -coloring of $E(G)$ has monochromatic copy of K_k
- $\mathcal{M}(k, r) :=$ all graphs G that are r -Ramsey minimal for K_k
(i.e., $G \rightarrow (K_k)_r$ and $G' \not\rightarrow (K_k)_r$ for all $G' \subsetneq G$)

Classical Ramsey-Parameters

- $\min_{G \in \mathcal{M}(k, 2)} v(G) =$ Ramsey number $R(k)$
- $\min_{G \in \mathcal{M}(k, 2)} e(G) =$ Size Ramsey number $\hat{R}(k)$

Minimum degree version (Burr-Erdős-Lovász 1976)

- $\min_{G \in \mathcal{M}(k, r)} \delta(G) =: s(k, r)$

Previous Results

- $s(k, 2) = (k - 1)^2$: Burr-Erdős-Lovász (1976)
- $s(k, r) = \tilde{\Theta}_k(r^2)$: Fox-Grinshpun-Liebenau-Person-Szabó (2015)

Fox-Grinshpun-Liebenau-Person-Szabó (JCTB, 2015)

$$cr^2 \log r \leq s(3, r) \leq Cr^2(\log r)^2$$

- They optimized their $s(k, r) = \tilde{\Theta}_k(r^2)$ bounds for $k = 3$
- Upper bound: they pack G_i sequentially via LLL-argument

Conjecture (Fox-Grinshpun-Liebenau-Person-Szabó, 2015)

$$s(3, r) = O(r^2 \log r)$$

- Hope: maybe can pack G_i sequentially via Δ -free process?
- For technical reasons (to make error-analysis tractable):
We pack G_i sequentially via semi-random Δ -free process

G., Warnke (2020): our packing result implies

$$s(3, r) = \Theta(r^2 \log r)$$